

L24 [$G = G_{n,p}$]

Technical Lemma (for pf. of Łuczak; source of 5/6)

$\exists \delta = \delta_n > 0$ s.t. whp

$M \subseteq V, |M| \leq n^{1/2 + \delta} =: m \implies$

(i) $|G[M]| < (3/2 - \delta) |M|$

(ii) $\chi(G[M]) \leq 3$

⏟
A

Pf indication

► (i) \rightarrow (ii) std obs (≠ culture):

$\chi(H) \leq \max_{K \subseteq H} \{ \chi(K) + 1 \} =: \text{col}(H)$ ("coloring #")

Pf = EX (neat, easy once found)

Pf of (i) = EX (U bd; AS L7.3.4)

NICE IDEA (Alon (?), improving SS):

$\varepsilon = \varepsilon(n) := 1/\log n$ [lots of room, need $\varepsilon \rightarrow 0$ fairly slowly]

► $k = k(n) = \min \{ l : \mathbb{P}(\chi(G) \leq l) \geq \varepsilon \}$

$\implies \chi(G) \geq k$ whp.

SHOW: $\chi(G) \leq \begin{cases} k+3 \\ k+1 \end{cases}$ whp

► $\chi = \chi(G) = \min \{ |S| : S \subseteq V, \chi(G-S) \leq k \}$ \rightarrow

$\mu = \mathbb{E}\chi$

• $\mathbb{P}(X=0) \geq \varepsilon$

• X Lip(1) (w.r.t. vert. exp.) $\xrightarrow{A_2} \forall \lambda > 0$

$$\left. \begin{array}{l} \mathbb{P}(X-\mu > \lambda) \\ \mathbb{P}(X-\mu < -\lambda) \end{array} \right\} < \exp[-\lambda^2/2n] \quad \text{or } n-1 \implies$$

$\varepsilon \leq \mathbb{P}(X=0) \leq e^{-\mu^2/2n} \implies$

• $\mu \leq \sqrt{2n \log_2(1/\varepsilon)} =: \lambda$

• whp $X \leq \mu + \lambda \leq 2\lambda$, i.e.

($\because \mathbb{P}(X \geq \mu + \lambda) \leq \exp[-\lambda^2/(2n)] = \varepsilon$)

whp $\exists S \subseteq V, |S| \leq 2\lambda$ s.t. $\chi(G-S) \leq k$ \textcircled{B}

\longrightarrow RTS Claim 1 $\textcircled{A} + \textcircled{B} \implies \chi(G) \leq k+3$

Claim 1 $\textcircled{A} + \textcircled{B} \implies \chi(G) \leq k+1$

Remark this is now deterministic

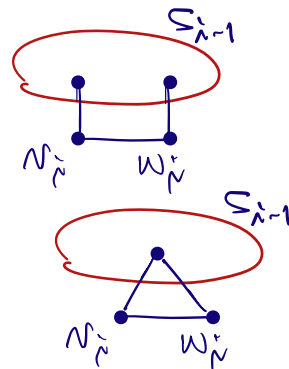
Pf of Cl. 1. S as in $\textcircled{B} \xrightarrow{\textcircled{A}} \chi(G) \leq k + \chi(G[S]) \leq k+3$ \textcircled{D}

Pf of Cl. 3 $S_0 = S$ (as in \textcircled{B}) \nexists do for $i \geq 1$ (until can't):

$S_i = S_{i-1} \cup \{v_i, w_i\} \cup \bar{v}$

$v_i, w_i \in V \setminus S_{i-1}, v_i \sim w_i \nexists v_i, w_i \sim S_{i-1}$

end at S_b



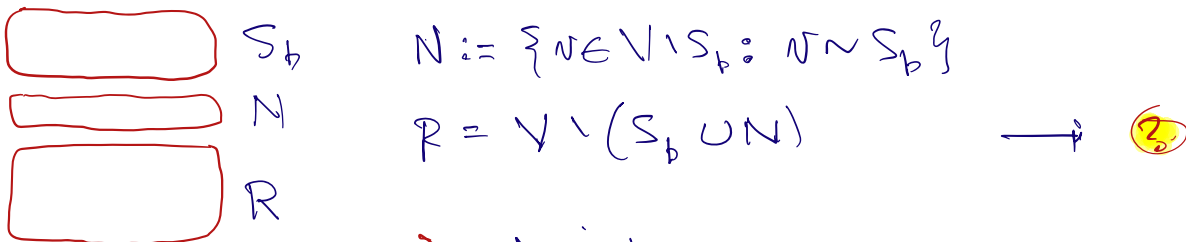
Obs $b < b_0 := \sqrt{n} \log n$ (want $n^{1/2+o(1)} > b_0 \gg \lambda$)

rough: each step adds 2 v's, 3 e's \rightarrow
if continue too long, violate $\textcircled{A}(i)$

precise: if reach b_0 then $|S_{b_0}| \leq 2\lambda + 2b_0$ ($< m$)

$$\nexists |G[S_{b_0}]| \geq 3b_0 > \left(\frac{\varepsilon}{2} - o(1)\right) |S_{b_0}| \quad \text{via } \textcircled{A}(i) \quad \square$$

End: $\chi(G[S_b]) \leq 3$ (by $\textcircled{A}(i)$)



N indept \rightarrow color:

R by $[k]$, N by $k+1$, S_b by $[3]$ \square

Small glitch: what if $k=2$? ($k=1$ silly)

Sketch/Ex: $k=2 \Rightarrow \mathbb{P}(\nexists \text{ odd cycle}) \geq \varepsilon \Rightarrow$
 $p < (1+o(1))/n \Rightarrow \text{col}(G) \leq 3$ whp. \square

Back to $\chi(G = G_{n, 1/2})$

start w $\alpha(G)$ (men'd in lec. 6)

• $f(k) (= f_n(k)) = \binom{n}{k} 2^{-\binom{k}{2}} (= \mathbb{E} |\{\text{ind } k\text{-sets}\}|)$

• $k_0 = k_0(n) = \max \{k : f(k) \geq 1\}$
 $= 2 \log n - 2 \log \log n + O(1)$

$\log = \log_2$

• $k = k_0 \pm O(1) \xrightarrow{\text{"Ex"}} \frac{f(k+1)}{f(k)} \approx \frac{k}{n} (= \frac{\log n}{n})$

• 2nd mm \rightarrow

a.s. $\alpha(G) \in \{k_0 - 1, k_0, k_0 + 1\}$

[\exists more: AS Cor 4.5.2]

\implies a.s. $\chi(G) > \frac{n}{2 \log n}$

\triangleright Balogh's 88: $\chi(G) \sim \frac{n}{2 \log n}$ a.s.

[i.e. $\forall \varepsilon > 0 \quad \mathbb{P}(\chi \neq (1 \pm \varepsilon) \frac{n}{2 \log n}) \rightarrow 0$]

naive:

$m = \lfloor \frac{n}{\log^2 n} \rfloor$

[just need: ① $m \ll n / \log^2 n$
② $\log m \sim \log n$]

• iterate: whp. $\exists I \in \mathcal{D}(G)$, $|I| \sim 2 \log n$

\rightarrow assign (new) color & delete

• STOP when $|V| \leq m$ ($|V| > m \implies \log |V| \sim \log n$)

& color rest w m new colors \rightarrow

$\chi(G) < (1 + o(1)) \frac{n}{2 \log n} + m \sim \frac{n}{2 \log n}$

$\square ?$

- NO! (why not?)

($\because G - I \neq G_{[n] \setminus I, 1/2}$)

▶ OTH: greedy ind. sets (equiv: gr. coloring) \rightarrow

$$\chi(G) < (1 + o(1)) \frac{n}{\log n} \quad \text{whp.}$$

[Ex to see why shd work — why is it different?]

MP for Ball: $\mathbb{P}(\alpha(G) < k_0 - \epsilon) = \exp[-\Omega(n^2 / \log^k n)]$

▶ Janson Ineq's ($\sim 90^\circ$ mot'd by Ball. but gen'l machine):

$$\log^8 \leftarrow \log^4$$

▶ conj \log^2 (EX: wd be tight) } irrelevant for Ball.

▶ then easy via "naive": (m as before)

MP $\xRightarrow{\otimes}$ whp. $\forall M \in \binom{[n]}{m}$

$$\alpha(G[M]) > \underbrace{k_0(m) - 3}_{\sim 2 \log n} \quad (m \geq 1)$$

$$\otimes \mathbb{P}(\text{tail}) < \binom{n}{m} \exp[-\Omega\left(\frac{m^2}{\log^k m}\right)] \quad (\frac{1}{2} \text{ no contest})$$

• Remk: role of MP analogous to that of "Tech. Lemma" in Łuczak