## Random matchings in regular graphs

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#### ABSTRACT

For a simple d-regular graph  $G$ , let M be chosen uniformly at random from the set of all matchings of G, and for  $x \in V(G)$  let  $p(\overline{x})$  be the probability that M does not cover x.

We show that for large d, the  $p(\bar{x})$ 's and the mean  $\mu$  and variance  $\sigma^2$  of  $|M|$  are determined to within small tolerances just by d and (in the case of  $\mu$  and  $\sigma^2$ )  $|V(G)|$ :

Theorem For any d-regular graph G,

(a)  $p(\overline{x}) \sim d^{-1/2}$   $\forall x \in V(G)$ , so that  $|V(G)| - 2\mu \sim |V(G)|/\sqrt{2}$ d,

(b)  $\sigma^2 \sim |V(G)|/(4\sqrt{d}),$ 

where the rates of convergence depend only on d.

## 1 Introduction

Given a graph  $G = (V, E)$ , write  $\mathcal{M}(G)$  for the set of matchings of G, and let M be chosen uniformly at random from  $\mathcal{M}(G)$ . (For graph theory background see e.g. [23]. We use "graph" to mean *simple* graph.) In this paper we are concerned with the behavior of  $M$ , and in particular of the random variable  $\xi = \xi_G = |M|$ , when G is regular of large degree.

Set  $p_k = p_k(G) = \Pr(\xi = k)$ . The distribution  $\{p_k\}$  (for a general G) has been considered in many contexts, in physics and chemistry as well as mathematics. We will not try to give a thorough bibliography, but see e.g. [19], [12], [22], [6], [7], [8], [23, Chapter 8].

These distributions are in some ways very nice. For instance, as shown in  $[11]$ ,  $[12]$ ,  $[22]$ , for any G the probability generating function

$$
f(G; \lambda) = \sum_{k} p_k \lambda^k \tag{1}
$$

has real roots. This gives log-concavity of the sequence  $\{p_k\}$  (*c.f.* "Newton's inequalities," e.g. [9, p.51]), and implies that the distribution is approximately normal provided the variance  $\sigma^2 = \sigma_{\xi}^2 =$  $\sigma^2(G)$  is large. (The latter is essentially due to L. Harper [10]. See the two paragraphs preceding Theorem 1.2 for some discussion and references concerning the question of when  $\sigma^2$  is large.)

Here we show that for regular G the behavior of  ${p_k}$  is nice in another sense: the mean  $(\mu = \mu_{\xi} =$  $\mu(G)$ ) and variance of  $\xi$  are remarkably well determined just by the degree and number of vertices of G.

Before stating this we need a finer parameter than  $\mu$ . For  $x \in V$ , write  $x \prec M$  if x is covered by (i.e. is contained in some edge of) the matching  $M$ , and set

$$
p(\overline{x}) = p_G(\overline{x}) = \Pr(x \nless M).
$$

Thus  $\mu = (n - \sum_{x \in V} p(\overline{x}))/2$ , where, here and throughout the paper, we set  $|V| = n$ .

**Theorem 1.1** For any d-regular graph  $G$ , (a)  $p(\overline{x}) \sim d^{-1/2}$   $\forall x \in V(G)$ , so that  $n - 2\mu(G) \sim n/\sqrt{d}$ , (b)  $\sigma^2(G) \sim n/(4\sqrt{d}).$ 

Here the limits are taken as  $d \to \infty$ ; so for example  $p(\bar{x}) \sim d^{-1/2}$  means

$$
(1 - o(1))d^{-1/2} < p(\overline{x}) < (1 + o(1))d^{-1/2},
$$

where  $o(1)$  depends only on d, and not on G or x. Let us stress that what's interesting here is the where  $o(1)$  depends only on a, and not on G or x. Let us stress that what s inter-<br>existence of the limiting values  $(d^{-1/2}, n/(4\sqrt{d}))$ , rather than the values themselves.

(The values themselves are easily seen to be a natural expression of the idea that the events  ${x \lt N}$  are roughly independent. To see this, we observe the easy identity (see (5))

$$
p(\overline{x}) = \left(1 + \sum_{y \sim x} p(\overline{y}|\overline{x})\right)^{-1}
$$

(where the conditional probability  $p(\bar{y}|\bar{x})$  has the obvious meaning). Using this, if we pretend the events  $\{x \prec M\}$  are mutually independent with  $p(\overline{x}) = p$  for all x, then

$$
p = (2d)^{-1}(-1 + \sqrt{1 + 4d}) = d^{-1/2} + (2d)^{-1} + O(d^{-3/2})
$$
\n(2)

gives (a) (see also Conjecture 1.3); while (b) derives from the fact that  $\xi$  is half the random variable  $|\{x \in V : x \prec M\}|$ , which has the binomial distribution  $B(n, 1-p)$ , so variance  $np(1-p) \sim nd^{-1/2}$ .

Let us also mention that it is not even easy to show that a large regular G has large  $\sigma^2(G)$ ; precisely: if  $G_{\alpha}$  is  $d_{\alpha}$ -regular  $(d_{\alpha} \neq 0)$  with  $n_{\alpha} := |V(G_{\alpha})| \to \infty$ , then  $\sigma^2(G_{\alpha}) \to \infty$  as  $\alpha \to \infty$ . This was shown in [6] provided  $d_{\alpha}/n_{\alpha} \to 0$ , but in full generality only in [17] (with a proof quite different from the arguments used here). So it is, again, rather surprising that one can say something as precise as Theorem 1.1.

(As more or less observed following (1), the condition  $\sigma^2(G_\alpha) \to \infty$  is equivalent to asymptotic normality of  $\{p_k(G_\alpha)\}_{k\geq 0}$ . This was the reason for most of the earlier work on  $\sigma^2(G)$ —see [27], [23, Ch. 8] in addition to [6], [17]—though not our principal motivation here.)

The actual bounds we establish are given in

**Theorem 1.2** For any d-regular graph G and  $\varepsilon > 0$ 

(a) 
$$
p(\overline{x}) = d^{-1/2} + O(d^{-3/4+\epsilon})
$$
  $\forall x \in V(G),$   
\n(b)  $\sigma^2(G) = (1 + O(d^{-1/4+\epsilon})) \frac{n}{4\sqrt{d}}$ .

Again note the error terms depend only on  $d$ . (This is slightly abusive, in a standard way. For example the error term in (a) does depend on G, x, but is bounded by some  $O(d^{-3/4+\epsilon})$  which depends only on d. So we should really write  $|p(\overline{x}) - d^{-1/2}| < O(d^{-3/4+\epsilon})$ .

Theorem 1.2 is proved beginning in Section 2. Before closing the present section, we just mention a few related questions.

First, it seems possible that the bounds in Theorem 1.2 can be strengthened considerably:

Conjecture 1.3 For any d-regular graph G and  $x \in V(G)$ , (a)  $p(\overline{x}) = d^{-1/2} - (2d)^{-1} + O(d^{-3/2}),$ (b)  $\sigma^2(G) = \frac{1}{4}nd^{-1/2} + O(nd^{-1}).$ 

(The coefficient of  $nd^{-1}$  in (b) is not an invariant.) This may be wishful thinking. It does, admittedly, seem too good to be true, but the same might have been (and was) said of Theorem 1.1 when it was not yet a theorem.

Second, it would be of considerable interest if something like Theorem 1.1 were true for hypergraphs of fixed edge size. We recall a few definitions. (For further background see e.g. [5] or [15].) A hypergraph  $H$  on a vertex set V is simply a collection of subsets of V, and is k-uniform if each of its members (called *edges*) is of size k. (So a 2-uniform hypergraph is a graph.) A hypergraph is d-regular if each of its vertices is contained in exactly  $d$  edges, and *simple* if no two of its vertices are contained in two distinct edges.

A matching in a hypergraph is again a collection of pairwise disjoint edges, and, as for graphs, we write  $\nu(\mathcal{H})$  for the size of a largest matching of H. We extend our earlier notation  $(\mathcal{M}, \xi, p(\overline{x}), \ldots)$  in the natural ways.

**Conjecture 1.4** Fix k. If  $H$  is a simple, k-uniform, d-regular hypergraph on a vertex set V of size n, then

(a) 
$$
p(\overline{x}) \sim d^{-1/k}
$$
  $\forall x \in V$ , so that  $n - k\mu(\mathcal{H}) \sim nd^{-1/k}$ , and in particular  $\mu(\mathcal{H}) \sim n/k$ ,  
(b)  $\sigma^2(\mathcal{H}) \sim n/(k^2 d^{1/k})$ 

(where again limits are taken as  $d \to \infty$ ).

This would be extremely interesting, not only for its own sake, but also because of its relation to work done over the last fifteen or so years on the asymptotic behavior of hypergraphs of bounded edge size. A central result in this area, proved by N. Pippenger following ideas of Ajtai, Komlós and Szemerédi  $[1]$ , Rödl  $[26]$  and Frankl and Rödl  $[4]$ , says in part:

**Theorem 1.5 (unpublished; see [28], [5])** Fix k. If  $H$  is as in Conjecture 1.4, then

$$
\nu(\mathcal{H}) > (1 - o(1))n/k,\tag{3}
$$

where  $o(1)$  depends only on d.

(See also e.g. [25], [5], [15], [14], [20], [18],[13], [29],[16], [21] for exposition and related work. For Pippenger's Theorem in full we should relax "simple" to a (uniform) bound  $o(d)$  on the pairwise degrees  $d(x, y)$ ; but we are in deep enough waters with the present hypotheses and will not explore this extra generality.)

Pippenger's Theorem is proved by a semirandom procedure (again, with roots in [1], [26]), and is not at all easy. So it would be extremely interesting if, as implied by Conjecture 1.4 (a), even  $\mu(\mathcal{H})$ could be shown to observe the behavior predicted for  $\nu(\mathcal{H})$  in (3). (This is of course much weaker than Conjecture 1.4 (a). In fact the error term  $(1+o(1))d^{-1/k}n/k$  (in the approximation  $\mu \approx n/k$ ) implied by the conjecture is far better than what one can at present establish in (3).)

In contrast, for a *graph* G as in Theorem 1.1: (i) Vizing's Theorem ([30] or e.g. [23, Theorem 7.4.1) implies  $\nu(G) \ge (1 - 1/(d+1))n/2$ ; and (ii)  $\mu > (1 - O(d^{-1/2}))n/2$ —a less precise version of Theorem 1.1 (a)—is not too hard to prove using the approach of Section 2 (see  $(8)$ ).

## 2 Path-trees and indication of proof

We first recall Godsil's [7] notion of the path-tree  $T(G, v)$  associated with a graph G and  $v \in V(G)$ . (This is called a *tree of walks* in [7]. The present name is from [23]. Were it not for its length, we would prefer "tree of self-avoiding walks," since we will eventually view the vertices of  $T(G, v)$  as outcomes of a random self-avoiding walk in  $G$ .)

The vertices of  $T = T(G, v)$  are the paths of G which begin at v. (For our purposes a path is a sequence  $(y_0, y_1, \ldots, y_l)$  of distinct vertices with  $y_i \sim y_{i-1}$ .) Two vertices of T are adjacent if one is a maximal proper subpath of the other.

We will usually use  $X, Y, Z, \ldots$  for vertices of T, and in particular write V for the singleton path  $(v)$ , which we regard as the root of T. Later we will be interested in a random path  $(v = y_0, y_1, \ldots, y_k)$ in G, and will write  $Y_l$  for the vertex  $(y_0, \ldots, y_l)$  (where  $l \leq k$ ). For  $W \in V(T)$  we write  $|W|$  for the length of the path W, in other words the depth of W in T, and set  $T_l = \{W \in V(T) : |W| = l\}$ . We use  $S(w)$  for the set of children of W,  $s(w)$  for  $|S(w)|$  and  $T(w)$  for the subtree rooted at W.

Path-trees  $T(G, v)$  turn out to capture considerable information about matchings in  $G$ , and to be in some respects easier to work with than the graph itself. (Again see [7] or the exposition in [23].) For present purposes the relevant connection is given by

**Lemma 2.1** With notation as above,  $p_G(\overline{v}) = p_{T(G,v)}(\overline{v})$ .

That is, the probability that a random matching of  $G$  misses  $v$  is the same as the probability that a random matching of T misses V.

Proof. This is an immediate consequence of the main result of [7], which we repeat here for the reader's convenience.

The matching generating polynomial of G is

$$
g(G;\lambda)=|\mathcal{M}(G)|f(G;\lambda)=\sum_k m_k\lambda^k
$$

where  $m_k = m_k(G)$  is the number of matchings of size k in G. The main result of [7] is (equivalent to)

$$
g(G - v; \lambda) / g(G; \lambda) = g(T - v; \lambda) / g(T; \lambda).
$$

Evaluation at  $\lambda = 1$  gives the lemma.

 $\blacksquare$ 

An advantage of working with  $T(G, v)$  is that it allows us to compute probabilities  $p_G(\overline{x})$  recursively. Let us extend our earlier notation, writing  $p(\overline{y}|\overline{x})$  for  $p_{G-x}(\overline{y})$  and  $p(\overline{x}, \overline{y})$  for  $Pr(x \nless M, y \nless M)$ . Since  $p(\overline{x}, \overline{y}) = p({x, y} \in M)$  when  $x \sim y$ , we have

$$
p(\overline{x}) + \sum_{y \sim x} p(\overline{x}, \overline{y}) = 1,\tag{4}
$$

which, when divided by  $p(\bar{x})$ , gives the basic identity

$$
p(\overline{x}) = \left(1 + \sum_{y \sim x} p(\overline{y}|\overline{x})\right)^{-1}.
$$
 (5)

For trees this takes the form

$$
p_{T(Y)}(\overline{Y}) = \left(1 + \sum_{Z \in S(Y)} p_{T(z)}(\overline{Z})\right)^{-1} \tag{6}
$$

where we write  $T(Y)$  for the subtree rooted at Y. Thus in principle we may compute the probabilities  $p_{T(Y)}(\overline{Y})$  recursively, beginning at the leaves and working up to the root, V, for which  $p_{T(Y)}(\overline{V}) = p_T(\overline{V})$ .

For example, if  $T = T(G, v)$  with G d-regular, then it's not hard to use this recursion together with the obvious

$$
d - l \le s(w) \le d \quad \forall w \in T_l \tag{7}
$$

to show

$$
c_1 d^{-1/2} < p_G(\overline{v}) < c_2 d^{-1/2} \quad \forall v \in V(G) \tag{8}
$$

for some positive constants  $c_1, c_2$ . (This gives the bound  $\mu(G) > (1 - O(d^{-1/2}))n/2$  mentioned at the end of Section 1.)

For Theorem 1.2 the inequalities (7) are not enough—e.g. the reader could try evaluating the extreme case

$$
s(\mathbf{w}) = \begin{cases} d & \text{if } |\mathbf{w}| \text{ is even} \\ d - |\mathbf{w}| & \text{if } |\mathbf{w}| \text{ is odd} \end{cases}
$$
(9)

—and we must show that degree fluctuations in  $T(G, v)$  are, in some usable sense, much more moderate than those in (9).

This is accomplished by comparing the degree  $s(W)$  of a vertex W with the average of the degrees of its children,

$$
\overline{s}(w) = \frac{1}{s(w)} \sum_{U \in S(w)} s(U).
$$

We show that, in contrast to (9),  $s(w)$  and  $\overline{s}(w)$  are close for most  $w \in V(T)$ . For the precise technical statement, set

$$
\Gamma(l,\varepsilon) := \{ w \in T_l : |s(w) - \overline{s}(w)| > d^{1/4+\varepsilon} \}
$$
 and  $\gamma(l,\varepsilon) := |\Gamma(l,\varepsilon)|$ ,

and let  $t = 4$  $\sqrt{d} \log d$ . (To prove (8) it's enough to consider something like the first  $\sqrt{d} \log d$  levels of T, and this will again be true for the proof of Theorem 1.2.)

**Lemma 2.2** For any fixed  $\varepsilon > 0$ , if d is sufficiently large and  $l \leq t$ , then

$$
\gamma(l,\varepsilon) < t^{-1}(d-t)^l e^{-d^{\varepsilon}}.\tag{10}
$$

This is proved in Section 3, and the derivation of Theorem 1.2 is completed in Sections 4 and 5. The bound (10) is given in a form convenient for later calculations, and is slightly weaker than what's produced in Section 3.

### 3 Proof of Lemma 2.2

For  $w \in V(G)$ , we write  $N(w)$  for the set of neighbors of w.

For  $w = (v = w_0, ..., w_l) \in V(T)$  let

$$
\delta(\mathbf{w}) = d - s(\mathbf{w}) = |\{w_0, \dots, w_{l-1}\} \cap N(w_l)|.
$$

Set

$$
\overline{\delta}(w) = \frac{1}{s(w)} \sum_{U \in S(w)} \delta(U) = d - \overline{s}(w).
$$

Our proof of Lemma 2.2 is more naturally expressed in terms of these parameters, that is, with  $\gamma(l,\varepsilon)$ rewritten as

$$
\gamma(l,\varepsilon) = |\{w \in T_l : |\delta(w) - \overline{\delta}(w)| > d^{1/4 + \varepsilon}\}|.
$$

Let  $(v = y_0, y_1, \ldots, y_t)$  be the natural random self-avoiding walk given by  $y_0 = v$  and

$$
\Pr(y_i = w | y_0, \dots, y_{i-1}) = s(y_{i-1})^{-1} 1_{\{w \in S(y_{i-1})\}},\tag{11}
$$

where, in agreement with our notation for  $T$ ,

$$
S(y_{i-1}) = N(y_{i-1}) \setminus \{y_0, \ldots, y_{i-2}\}\
$$

and  $s(y_{i-1}) = |S(y_{i-1})|$ ; that is, the walk chooses  $y_i$  uniformly from the as yet unvisited neighbors of  $y_{i-1}$ .

As earlier, we write  $Y_l$  for  $(y_0, \ldots, y_l)$ , thought of as a random vertex of  $T_l$ . We will show that for  $Y_l$  chosen according to this (not quite uniform) distribution on  $T_l$ ,  $|\delta(Y_l) - \delta(Y_l)|$  is very unlikely to be large; precisely, for any  $\alpha > 0$ ,

$$
\Pr(|\delta(\mathbf{Y}_l) - \overline{\delta}(\mathbf{Y}_l)| > \alpha + 32\log^2 d) < 2d^2 t \exp\left(-\frac{\alpha^2}{2t}\right). \tag{12}
$$

To see that this implies Lemma 2.2, note that for any  $w \in T_l$ ,

$$
\Pr(\mathbf{Y}_l = \mathbf{W}) \ge d^{-1}(d-1)^{-(l-1)} > d^{-l}
$$

whence, setting

$$
W_{\alpha} = \{ w \in T_l : |\delta(w) - \overline{\delta}(w)| > \alpha + 32 \log^2 d \},
$$

we have

$$
|W_{\alpha}| \le \Pr(\mathbf{Y}_l \in W_{\alpha}) \left( \min_{\mathbf{W} \in T_l} \Pr(\mathbf{Y}_l = \mathbf{W}) \right)^{-1} < 2d^2 t \exp\left(-\frac{\alpha^2}{2t}\right) d^l.
$$

Taking  $\alpha = d^{1/4+\epsilon} - 32 \log^2 d$  then gives Lemma 2.2 (and a bit more).

The key observation for the proof of (12) is that while  $\delta(Y_i)$  is the number of visits to  $N(y_i)$  by  $(y_0, \ldots, y_{l-1}), \delta(Y_l)$  is roughly the "expected" number of such visits, where "expected" is used in the dynamic sense given by the function  $f$  below. A little martingale analysis then shows that these actual and expected numbers are likely to be close.

For fixed  $l \in [t]$  and  $w \in V(G)$ , define

$$
f(w) = \sum_{i=1}^{l} \Pr(y_i \in N(w) | y_0, \dots, y_{i-1}),
$$
\n(13)

$$
g(w) = |N(w) \cap \{y_1, \dots, y_l\}|. \tag{14}
$$

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#### **Lemma 3.1** For any  $\alpha > 0$

$$
\Pr(\exists w \text{ with } |f(w) - g(w)| > \alpha) < 2d^2 t \exp\left(-\frac{\alpha^2}{2t}\right). \tag{15}
$$

Remark. The reader may observe below that we only use the fact that  $|f(w) - g(w)|$  is usually small when  $w = y_i$ ; but the proof gives the stated inequality, and in fact we don't see how to establish what we need for  $y_l$  without proving something like (15).

Before proving Lemma 3.1, let us see why it implies (12). Notice that

$$
g(y_i) = \delta(\mathbf{Y}_i). \tag{16}
$$

On the other hand, we show that  $f(y_i)$  is a good approximation of  $\delta(Y_i)$ . We have

$$
\overline{\delta}(\mathbf{Y}_l) = \frac{1}{s(y_l)} \sum \{ |N(u) \cap \{y_0, \dots, y_l\}| : u \in N(y_l) \setminus \{y_0, \dots, y_{l-1}\} \},\tag{17}
$$

while a similar expression for  $f(w)$  is

$$
\sum_{i=1}^{l} \frac{1}{s(y_{i-1})} |(N(y_{i-1}) \cap N(w)) \setminus \{y_0, \dots, y_{i-2}\}|. \tag{18}
$$

Now when  $w = y_i$ , the sum of the set cardinalities appearing in (18) is not much different than the sum in (17): the former—that is,

$$
\sum_{i=1}^{l} |(N(y_{i-1}) \cap N(y_i)) \setminus \{y_0, \dots, y_{i-2}\}|
$$
\n(19)

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—counts ordered pairs  $(u, y_{i-1})$  with  $1 \leq i \leq l$ ,  $y_i \sim u \sim y_{i-1}$ , and  $u \notin \{y_0, \ldots, y_{i-2}\}$ ; whereas the latter counts all such pairs for which  $u \notin \{y_0, \ldots, y_{l-1}\}$ , together with the pairs  $(u, y_l)$  with  $u \in N(y_{_l}) \setminus \{y_0, \ldots, y_{_{l-1}}\}.$ 

The difference between these sums is thus bounded by

$$
\max \Big\{ |\{(j,i) : i \le j \le l-1, y_l \sim y_j \sim y_{i-1} \}|, d \Big\} \le \binom{t}{2},
$$

and we have (using (7))

$$
|f(y_l) - \overline{\delta}(\mathbf{Y}_l)| \leq \frac{1}{s(y_l)} \binom{t}{2} + \sum_{i=1}^l \left| \frac{1}{s(y_{i-1})} - \frac{1}{s(y_l)} \right| |N(y_{i-1}) \cap N(y_l)|
$$
  

$$
\leq \frac{1}{d-t} \binom{t}{2} + \left( \frac{1}{d-t} - \frac{1}{d} \right) ld
$$
  

$$
< 2t^2 d^{-1} \leq 32 \log^2 d.
$$

Of course this together with (16) shows that Lemma 3.1 implies (12).

*Proof of Lemma 3.1.* Let us for the moment fix  $w \in V$  and write

$$
f(w) - g(w) = \sum_{i=1}^{l} X_i
$$

where

$$
X_i = X_i(w) = \Pr(y_i \in N(w)|y_0, \dots, y_{i-1}) - 1_{\{y_i \in N(w)\}}.
$$

Now  $\{X_i\}_{i=1}^l$  is a martingale difference sequence (that is,  $E[X_i|X_1,\ldots,X_{i-1}]=0$ ), with

$$
|X_i| \le 1. \tag{20}
$$

So according to "Azuma's inequality" (see, e.g., [3], [24], [2]), for any  $\alpha > 0$ ,

$$
\Pr(|f(w) - g(w)| > \alpha) < 2 \exp\left(-\frac{\alpha^2}{2t}\right). \tag{21}
$$

Thus we have a bound like  $(15)$  for any fixed w.

For  $(15)$  we must somehow control the number of w's under consideration. A priori this number could be something like the number of vertices within distance t of  $v$  (which swamps the bound in (21)); but we can reduce it by only beginning to keep track of  $f(w) - g(w)$  when (and if) our random walk gets to within distance 2 of  $w$ .

To do this, let us fix, solely for bookkeeping purposes, some linear ordering "≺" of V . For each  $w \in V$  define the random variable  $i(w)$  by

$$
j(w) = \begin{cases} 0 & \text{if } d(v, w) \le 2\\ \infty & \text{if } d(y_i, w) > 2\\ \min\{i : d(y_i, w) = 2\} & \text{otherwise,} \end{cases}
$$

and then let  $v_s$  be the  $s^{\text{th}}$  vertex in the (lexicographic) ordering in which w precedes w' if either  $j(w) < j(w')$  or  $j(w) = j(w')$  and  $w \prec w'$ . (Note this is a random ordering determined by  $(y_0, \ldots, y_t)$ .) Now for  $1 \leq s \leq d^2t$  and  $1 \leq i \leq t$ , set

$$
X_i^s = X_i(v_s).
$$

(Note  $X_i^s = 0$  if  $i \leq j(v_s)$ . We could omit the restriction  $s \leq d^2t$ , but this adds nothing since for larger s we have  $j(v_s) = \infty$  and so  $X_i^s = 0$  for all *i*.)

Now for each fixed s,  $f(v_s) - g(v_s) = \sum_{i=1}^{l} X_i^s$ , and  $\{X_i^s\}_{i=1}^{l}$  is again a martingale difference sequence satisfying (20). Thus

$$
Pr(|f(v_s) - g(v_s)| > \alpha) < 2 \exp\left(-\frac{\alpha^2}{2t}\right)
$$

for each s, and

$$
\Pr(\exists s \in [d^2t], \ |f(v_s) - g(v_s)| > \alpha) < 2d^2t \exp\left(-\frac{\alpha^2}{2t}\right).
$$

But this gives (15), since (trivially)  $f(v_s) = g(v_s) = 0$  if  $s > d^2t$ .

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# 4 Proof of Theorem 1.2(a)

Fix  $v \in V(G)$  and set  $T := T(G, v)$  as in the preceding section. Let  $0 < \varepsilon < 0.1$  (fixed), and

$$
\eta := d^{1/4 + \varepsilon} , \quad \Gamma_i := \Gamma(i, \varepsilon) \quad i = 0, 1, ..., t .
$$

Then Lemma 2.2 says that

$$
|\Gamma_i| < t^{-1}(d-t)^i e^{-d^\varepsilon} \tag{22}
$$

Define, for  $0 \leq i \leq j \leq t$ ,

$$
B_{i,j} := \{ w \in T_i : |T(w) \cap \Gamma_j| \ge t^{-1} (d-t)^{j-i} e^{-d^{\varepsilon}/2} \} \text{ and } B_i := \bigcup_{j=i}^t B_{i,j} .
$$
 (23)

Then (22) yields  $B_0 = \emptyset$  and for  $i \geq 1$ ,

$$
|B_i| \le \sum_{j=i}^t |B_{i,j}| \le \sum_{j=i}^t \frac{|\Gamma_j|}{t^{-1}(d-t)^{j-i}e^{-d^{\varepsilon}/2}} \le t(d-t)^i e^{-d^{\varepsilon}/2} . \tag{24}
$$

Let  $q_T(\overline{w}) = p_{T(w)}(\overline{w})$ . Set  $q_t = q_{t-1} = 1$  and for  $i = 0, 1, \ldots, t-2$ ,

$$
q_i = \left(1 + \frac{d-i}{1 + (d-i+\eta)q_{i+2}}\right)^{-1}.
$$

We claim

$$
q_i = \frac{1}{\sqrt{d}} + O(d^{-3/4 + \varepsilon}) \quad \text{for } 0 \le i \le t/2 + 1. \tag{25}
$$

**Proof.** We prove this for i even; odd i is handled similarly. For  $i = 0, 2, \dots, t$ , define

$$
f_i(x) := \left(1 + \frac{d-i}{1 + (d-i+\eta)x}\right)^{-1}, \quad x > 0
$$

(so  $f_i(q_{i+2}) = q_i$ ) and denote by  $a_i$  the unique positive solution of  $f_i(a_i) = a_i$ ; that is,

$$
a_i := \frac{\eta - 1 + \sqrt{(\eta - 1)^2 + 4(d - i + \eta)}}{2(d - i + \eta)}.
$$

It is easy to check that

$$
a_i = \frac{1}{\sqrt{d}} + O(d^{-3/4+\epsilon})
$$
 and  $0 < a_{i+2} - a_i \le d^{-3/2}$ .

Thus it is enough to show by reverse induction that

$$
0 \le q_i - a_i \le \exp\left(-\frac{3(t-i)}{4\sqrt{d}}\right) + \frac{(t-i)d^{-3/2}}{2} \,,\tag{26}
$$

for  $i = t, t - 2, \dots, 0$ . (Recall that  $t = 4$ ) √  $d \log d$ . )

The base case  $i = t$  is trivial. Suppose (26) is true for  $i + 2 \leq t$ . Since  $f_i$  is increasing, the lower bound of the induction hypothesis gives

$$
a_i = f_i(a_i) \le f_i(a_{i+2}) \le f_i(q_{i+2}) = q_i.
$$

On the other hand, the Mean Value Theorem implies that there exists x with  $a_i \leq x \leq q_{i+2}$  and such that

$$
q_i - a_i = f_i(q_{i+2}) - f_i(a_i) = f'_i(x)(q_{i+2} - a_i) .
$$

Since  $f_i'(z) \leq e^{-3/(2\sqrt{d})}$  for  $z \geq a_i = 1/$  $\overline{d} + O(d^{-3/4+\epsilon})$ , we have

$$
q_i - a_i \le e^{-3/(2\sqrt{d})} \left( \exp\left(-\frac{3(t-i-2)}{4\sqrt{d}}\right) + \frac{(t-i-2)d^{-3/2}}{2} + d^{-3/2} \right)
$$
  

$$
\le \exp\left(-\frac{3(t-i)}{4\sqrt{d}}\right) + \frac{(t-i)d^{-3/2}}{2}.
$$

The following lemma and its corollary (which includes Theorem 1.2(a)) will again be needed in the proof of Theorem 1.2(b).

**Lemma 4.1** Let  $0 \leq i \leq t-1$  and  $x \in T_i \setminus B_i$ . Then

$$
q_T(\overline{\mathbf{x}}) \le q_i + e^{-d^{\varepsilon}/2} \tag{27}
$$

 $\Box$ 

and

$$
q_T(\overline{\mathbf{w}}) \le q_{i+1} + (d-t)e^{-d^{\varepsilon}/2} \quad \text{for all } \mathbf{w} \in S(\mathbf{x}).\tag{28}
$$

Note that (6) and (28) yield

$$
q_T(\overline{\mathbf{x}}) \ge \left(1 + d(q_{i+1} + (d-t)e^{-d^{\varepsilon}/2})\right)^{-1} \quad \text{for } \mathbf{x} \in T_i \setminus B_i \, .
$$

Thus the following corollary follows from Lemma 4.1 and (25).

**Corollary 4.2** Let  $0 \leq i \leq t/2$  and  $x \in T_i \setminus B_i$ . Then

$$
\left| q_{T}(\overline{\mathbf{x}}) - \frac{1}{\sqrt{d}} \right| = O(d^{-3/4+\varepsilon}).
$$

In particular, since  $B_0 = \emptyset$  and  $q_T(\overline{v}) = p_T(\overline{v})$  we have Theorem 1.2(a).

We prove (27) for |x| even (that is, congruent mod 2 to t) and (28) for |x| odd (that is, for |w| even). The proof when these parities are reversed is identical, except that one truncates at odd rather than even levels in the following definition of  $T'$ .

We first define an auxiliary subtree  $T'$  of T. Let

$$
\Gamma := \bigcup_{i \text{ even}} \Gamma_i \cup T_t.
$$

The auxiliary tree  $T'$  is the tree obtained from T by removing all descendants of vertices in  $\Gamma$ . Note that all leaves of T' are in  $\Gamma$  (T itself has no leaves above level d) and that  $X \in T'$  (since  $X \notin B_{|X|}$ implies  $x \notin \Gamma$ ). For  $w \in V(T')$ , let

$$
q(\overline{\mathbf{w}}) = q_T(\overline{\mathbf{w}}) \quad \text{and} \quad q'(\overline{\mathbf{w}}) = q_{T'}(\overline{\mathbf{w}}).
$$

Since for each leaf w of T', |w| is even and  $q(\overline{w}) \leq q'(\overline{w}) = 1$ , it follows easily from (6) that

 $q(\overline{w}) \leq q'(\overline{w})$  for all  $w \in V(T')$  with  $|w|$  even.

In particular,

$$
q(\overline{\mathbf{x}}) \le q'(\overline{\mathbf{x}}) \tag{29}
$$

We prove the following lemma from which (27) follows easily.

**Lemma 4.3** For all  $w \in T'$  with  $|w|$  even, we have

$$
q'(\overline{\mathbf{w}}) \le q_{|\mathbf{w}|} + \sum_{\mathbf{v} \in L'(\mathbf{w})} (d-t)^{-|\mathbf{v}|+|\mathbf{w}|} (1-q_{|\mathbf{v}|}), \qquad (30)
$$

where  $L'(w)$  is the set of leaves of  $T'(w)$ .

(We regard the vertex of a singleton tree as a leaf.)

*Proof.* If w is a leaf of T', (30) is trivial (the sum in (30) is  $1 - q_{|W|}$ ). Suppose |W| is even and (30) is true for all descendants of W with even length. Then (6) gives, using Jensen's inequality,

$$
q'(\overline{w}) = \left(1 + \sum_{U \in S(W)} q'(\overline{U})\right)^{-1}
$$
  
= 
$$
\left(1 + \sum_{U \in S(W)} \frac{1}{1 + \sum_{X \in S(U)} q'(\overline{X})}\right)^{-1}
$$
  

$$
\leq \left(1 + \frac{s(W)}{1 + s(W)^{-1} \sum_{U \in S(W)} \sum_{X \in S(U)} q'(\overline{X})}\right)^{-1}.
$$

The induction hypothesis yields

$$
s(w)^{-1} \sum_{v \in S(w)} \sum_{x \in S(v)} q'(\overline{x})
$$
  
\n
$$
\leq s(w)^{-1} \sum_{v \in S(w)} \sum_{x \in S(v)} q_{|x|} + s(w)^{-1} \sum_{v \in S(w)} \sum_{x \in U} \sum_{x \in U(x)} (d-t)^{-|z|+|x|} (1 - q_{|z|})
$$
  
\n
$$
= \overline{s}(w) q_{|w|+2} + s(w)^{-1} \sum_{z \in L'(w)} (d-t)^{-|z|+|w|+2} (1 - q_{|z|})
$$
  
\n
$$
\leq (s(w) + \eta) q_{|w|+2} + \sum_{z \in L'(w)} (d-t)^{-|z|+|w|+1} (1 - q_{|z|}),
$$

where the last inequality uses W  $\notin \Gamma$  (since W is not a leaf) and  $s(w) \geq d - |w| \geq d - t$ .

Since

$$
\left(1 + \frac{\alpha}{\beta + x}\right)^{-1} \le \left(1 + \frac{\alpha}{\beta}\right)^{-1} + \alpha^{-1}x \text{ for all } \alpha, \beta, x > 0,
$$

we have

$$
q'(\overline{w}) \leq \left(1 + \frac{s(w)}{1 + (s(w) + \eta)q_{|W|+2} + \sum_{z \in L'(w)} (d-t)^{-|z|+|W|+1} (1 - q_{|z|})}\right)^{-1}
$$
  
\n
$$
\leq \left(1 + \frac{s(w)}{1 + (s(w) + \eta)q_{|W|+2}}\right)^{-1} + s(w)^{-1} \sum_{z \in L'(w)} (d-t)^{-|z|+|W|+1} (1 - q_{|z|})
$$
  
\n
$$
\leq \left(1 + \frac{d-|w|}{1 + (d-|w| + \eta)q_{|W|+2}}\right)^{-1} + \sum_{z \in L'(w)} (d-t)^{-|z|+|w|} (1 - q_{|z|})
$$
  
\n
$$
= q_{|w|} + \sum_{z \in L'(w)} (d-t)^{-|z|+|w|} (1 - q_{|z|}).
$$

Proof of (27). Because Lemma 4.3 and (29) give

$$
q(\overline{x}) \le q_{|x|} + \sum_{U \in L'(x)} (d-t)^{-|U|+|x|} (1-q_{|U|}),
$$

 $\blacksquare$ 

 $\blacksquare$ 

it is enough to show that

$$
\sum_{\mathbf{U}\in L'(\mathbf{X})}(d-t)^{-|\mathbf{U}|+|\mathbf{X}|}(1-q_{|\mathbf{U}|})\leq e^{-d^{\varepsilon}/2}\ .
$$

But  $X \notin B_{|X|}$  and  $q_t = 1$  imply that

$$
\sum_{\substack{U \in L'(x) \\ |x| \le j \le t-2}} (d-t)^{-|U|+|X|} (1-q_{|U|}) \le \sum_{\substack{j : \text{ even} \\ |x| \le j \le t-2}} \sum_{\substack{U \in T(x) \cap \Gamma_j \\ t^{-1} (d-t)^{j-|X|} e^{-d^{\varepsilon}/2} (d-t)^{-j+|X|} \le e^{-d^{\varepsilon}/2}}} \cdot (d-t)^{-j+|X|} \le e^{-d^{\varepsilon}/2}.
$$

The proof of  $(28)$  is the same as that of  $(27)$ , except we use

$$
|T(\mathbf{W}) \cap \Gamma_j| \le t^{-1}(d-t)^{j-|\mathbf{W}|+1}e^{-d^{\varepsilon}/2}
$$

(which follows from  $w \in S(x)$  and  $x \notin B_{|X|}$ ).

# 5 Proof of Theorem 1.2(b)

For a graph G and  $v \in V(G)$ , set

$$
I(G, v) = \sum_{w \in V(G) \setminus \{v\}} \left( p_G(\overline{w}|\overline{v}) - p_G(\overline{w}) \right) .
$$

Notice that

$$
\sigma^2(G) = \text{Var}\left[\frac{1}{2}\left(n - \sum_{v \in V(G)} 1_{\{v \neq M\}}\right)\right] = \frac{1}{4}\text{Var}\left[\sum_{v \in V(G)} 1_{\{v \neq M\}}\right].
$$

On the other hand, using Theorem  $1.2(a)$ , we have

$$
\operatorname{Var}\left[\sum_{v \in V(G)} 1_{\{v \neq M\}}\right] = \sum_{v,w \in V(G)} \left(p_G(\overline{v}, \overline{w}) - p_G(\overline{v})p_G(\overline{w})\right)
$$
  
\n
$$
\leq \sum_{v \in V(G)} p_G(\overline{v}) + \sum_{v \in V(G)} p_G(\overline{v}) \sum_{w \in V(G) \setminus \{v\}} \left(p_G(\overline{w}|\overline{v}) - p_G(\overline{w})\right)
$$
  
\n
$$
= \left(1 + O(d^{-1/4+\varepsilon})\right) \frac{n}{\sqrt{d}} + \sum_{v \in V(G)} p(\overline{v})I(G, v) . \tag{31}
$$

Thus Theorem 1.2(b) will follow from

$$
|p(\overline{v})I(G,v)| \le 2d^{-3/4+4\varepsilon} \qquad \text{for all } v \in V(G). \tag{32}
$$

Of course (31) and (32) are a concrete expression of the idea that the indicators  $1_{\{v \not\leq M\}}$  are close to independent (compare the discussion in the vicinity of (2) of the limiting values in Theorem 1.2).

We will use the following lemmas in the proof of (32).

**Lemma 5.1** For any graph  $G$ ,

$$
I(G,v) = -\sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) I(G_v, y) + \sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) p_G(\overline{y}|\overline{v}),
$$

where  $G_v$  is the subgraph of G induced by  $V(G) \setminus \{v\}.$ 

*Proof.* Notice that, for  $e_{y,v} = \{y, v\}$  and  $p_G(e_{y,v}) := \Pr_G(e_{y,v} \in M)$ ,

$$
p_G(\overline{v}) + \sum_{y \in N(G)} p_G(e_{y,v}) = 1,
$$

and

$$
p_G(\overline{x}) = p_G(\overline{v})p_G(\overline{x}|\overline{v}) + \sum_{y \in N(G)} p_G(e_{y,v})p_G(\overline{x}|e_{y,v} \in M) .
$$

Thus

$$
p_G(\overline{x}|\overline{v}) - p_G(\overline{x}) = p_G(\overline{x}|\overline{v}) \left( p_G(\overline{v}) + \sum_{y \in N_G(v)} p_G(e_{y,v}) \right) - p_G(\overline{x})
$$
  

$$
= - \sum_{y \in N_G(v)} p_G(e_{y,v}) \left( p_G(\overline{x}|e_{y,v} \in M) - p_G(\overline{x}|\overline{v}) \right).
$$

Since  $p_G(e_{y,v}) = p_G(\overline{y}, \overline{v})$  and

$$
p_G(\overline{x}|e_{y,v} \in M) = \begin{cases} p_G(\overline{x}|\overline{y}, \overline{v}) & \text{if } x \notin \{y, v\} \\ 0 & \text{if } x \in \{y, v\}, \end{cases}
$$

we have

$$
I(G, v) = -\sum_{x \in V(G) \backslash \{v\}} \sum_{y \in N_G(v)} p_G(e_{y,v}) \left( p_G(\overline{x}|e_{y,v} \in M) - p_G(\overline{x}|\overline{v}) \right)
$$
  
\n
$$
= -\sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) \sum_{x \in V(G) \backslash \{v\}} \left( p_G(\overline{x}|e_{y,v} \in M) - p_G(\overline{x}|\overline{v}) \right)
$$
  
\n
$$
= -\sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) \sum_{x \in V(G) \backslash \{y,v\}} \left( p_G(\overline{x}|\overline{y}, \overline{v}) - p_G(\overline{x}|\overline{v}) \right) + \sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) p_G(\overline{y}|\overline{v}) .
$$

Because

$$
p_G(\overline{x}|\overline{v}) = p_{G_v}(\overline{x})
$$
 and  $p_G(\overline{x}|\overline{y}, \overline{v}) = p_{G_v}(\overline{x}|\overline{y})$ ,

this yields

$$
I(G, v) = -\sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) I(G_v, y) + \sum_{y \in N_G(v)} p_G(\overline{y}, \overline{v}) p_G(\overline{y}|\overline{v}) .
$$

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For  $w = (v = y_0, y_1, ..., y_k) \in T := T(G, v)$ , let

$$
p_{\boldsymbol{G}}(\overline{\mathbf{W}}) = \mathbf{Pr}_{{\boldsymbol{G}}}(y_i \not\prec M \quad \forall \; i=0,...,k),
$$

and for  $0 \leq i \leq |w|$ , set  $W_i := (v = y_0, ..., y_i)$  (the ancestor of w having length i). Central to our argument are the quantities

$$
r(\overline{\mathbf{w}}) := \prod_{i=0}^{|w|} q(\overline{\mathbf{w}}_i) = p_G(\overline{\mathbf{w}}),
$$

where, as in Section 4, we write  $q(\overline{w})$  for  $p_{T(w)}(\overline{w})$ .

In what follows we often use unsubscripted p's (e.g.  $p(\overline{v})$ ,  $p(\overline{y}|\overline{v})$ ) for the corresponding  $p_G$ 's.

Lemma 5.2 With notation as above,

$$
p_G(\overline{v}) = \sum_{\mathbf{W} \in V(T)} r^2(\overline{\mathbf{W}}) .
$$

*Proof.* We induct on  $|V(G)|$ , the case  $|V(G)| = 1$  being trivial. Note that for each  $W = (v, w) \in S(V)$ , we have  $T(w) = T(G - v, w)$ , and that our induction hypothesis expressed in terms of  $T(w)$  is

$$
p_{T(\mathbf{W})}(\overline{\mathbf{W}}) = \sum_{\mathbf{X} \in V(T(\mathbf{W}))} \frac{r^2(\overline{\mathbf{X}})}{p_G^2(\overline{v})} .
$$

Thus Lemma 2.1 and (6) give

$$
p(\overline{v}) = p_T(\overline{v}) = [1 + \sum_{w \in S(v)} p_{T(w)}(\overline{w})]^{-1} = [1 + \sum_{x \in V(T) \setminus \{v\}} \frac{r^2(\overline{x})}{p^2(\overline{v})}]^{-1}
$$

which is the same as

$$
p(\overline{v}) = p^2(\overline{v}) + \sum_{\mathbf{x} \in V(T) \setminus \{v\}} r^2(\overline{\mathbf{x}})
$$

(which is what we want).

Lemma 5.3 With notation as above,

$$
p_G(\overline{v})I(G,v) = \sum_{w \in V(T) \backslash \{v\}} (-1)^{|w|-1} r^2(\overline{w}) .
$$

*Proof.* We again induct on  $|V(G)|$  (with the base case  $|V(G)| = 1$  again trivial).

Using Lemma 5.1 (and  $p(\overline{v}, \overline{y}) = p(\overline{v})p(\overline{y}|\overline{v})$ ) we have

$$
p(\overline{v})I(G,v) = -p^2(\overline{v})\sum_{y \in N_G(v)} p(\overline{y}|\overline{v})I(G_v,y) + \sum_{y \in N_G(v)} p^2(\overline{y},\overline{v}).
$$
\n(33)

The second sum is just

$$
\sum_{\substack{Y \in V(T) \\ |Y| = 1}} r^2(Y).
$$

For the first sum, note that for  $y \in N_G(v)$  and  $Y = (v, y) \in V(T)$ , our inductive hypothesis says that

$$
p(\overline{y}|\overline{v})I(G_v, y) = \sum_{\mathbf{x}\in V(T(\mathbf{Y}))\backslash{\{\mathbf{Y}\}}} (-1)^{|\mathbf{x}|} \frac{r^2(\overline{\mathbf{x}})}{p^2(\overline{v})}.
$$

(Here |x| and  $r(\bar{x})$  refer to T, so the correction  $-p^{-2}(\bar{v})$  gives the appropriate expression for  $T(y)$ .) Thus the first term on the right hand side of (33) is

$$
-\sum_{Y\in S(V)}\sum_{X\in V(T(Y))\setminus\{Y\}}(-1)^{|X|}r^2(\overline{x})=\sum_{\substack{X\in V(T)\\|X|\geq 2}}(-1)^{|X|-1}r^2(\overline{x}),
$$

so we have Lemma 5.3.

*Proof of Theorem* 1.2(b). We will show that, with  $p(\overline{v}) := p_T(\overline{v}) \ (= p_G(\overline{v})),$ 

$$
\sum_{\substack{w \in V(T) \setminus \{v\} \\ |w| \text{ even}}} r^2(\overline{w}) \ge \frac{p(\overline{v})(1 - d^{-1/4 + 4\varepsilon})}{2} \tag{34}
$$

and

$$
\sum_{\substack{W \in V(T) \\ |W| \text{ odd}}} r^2(\overline{W}) \ge \frac{p(\overline{V})(1 - d^{-1/4 + 4\varepsilon})}{2} \ . \tag{35}
$$

(We wind up with the  $(4\varepsilon)$ 's because we often use extra factors  $d^{\varepsilon}$  to subsume smaller but clumsier error terms.) The inequalities (34), (35), together with Theorem 1.2(a) and Lemma 5.2 imply that

$$
\sum_{\mathbf{w}\in V(T)\backslash\{\mathbf{v}\}} (-1)^{|\mathbf{w}|-1} r^2(\overline{\mathbf{w}}) = \sum_{\mathbf{w}\in V(T)\backslash\{\mathbf{v}\}} r^2(\overline{\mathbf{w}}) - 2 \sum_{\mathbf{w}\in V(T)\backslash\{\mathbf{v}\}} r^2(\overline{\mathbf{w}})
$$
\n
$$
\leq p(\overline{\mathbf{v}}) - 2 \cdot \frac{p(\overline{\mathbf{v}})(1 - d^{-1/4 + 4\varepsilon})}{2} \leq 2d^{-3/4 + 4\varepsilon}
$$

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and similarly

$$
\sum_{w \in V(T) \setminus \{v\}} (-1)^{|w|-1} r^2(\overline{w}) \ge -2d^{-3/4+4\varepsilon} ,
$$

which together with Lemma 5.3 yield (32).

The proofs of (34) and (35) are essentially identical, so we only prove (34). Set  $a(v) = 1$  and

$$
a(\mathbf{w}) = \left(\prod_{\substack{0 \le i \le l-2 \\ i \text{ even}}} s(\mathbf{w}_i)\right)^{-1} \quad \text{for } \mathbf{w} \in V(T) \text{ with } l := |\mathbf{w}| \ge 2 \text{ even.}
$$

(Recall  $s(w) = |S(w)|$ , the number of children of w.) The Cauchy-Schwarz inequality says

$$
\left(\sum_{w \in T_l} a(w)r(\overline{w})\right)^2 \le \sum_{w \in T_l} a^2(w) \sum_{w \in T_l} r^2(\overline{w}) . \tag{36}
$$

We will prove (34) by establishing, in Claims 1 and 2 below, lower and upper bounds on (respectively) the left hand side of (36) and the first term on the right hand side.

### Claim 1.

$$
\sum_{w \in T_l} a(w)r(\overline{w}) \ge p(\overline{v}) \left(1 - d^{-1/2} - 2d^{-3/4 + 2\varepsilon}\right)^{l/2} \quad \text{for even } l \le t/2 - 2.
$$

*Proof.* We show by induction (the case  $l = 0$  being trivial) that

$$
\sum_{\mathbf{W}\in T_l} a(\mathbf{W}) r(\overline{\mathbf{W}}) \ge p(\overline{\mathbf{V}}) \left(1 - d^{-1/2} - d^{-3/4+2\varepsilon}\right)^{l/2} - \frac{l \cdot e^{-d\varepsilon/5}}{2}.
$$

Note that (6) gives

$$
\sum_{w \in T_{l+2}} a(w)r(\overline{w}) = \sum_{x \in T_l} \sum_{v \in S(x)} \sum_{w \in S(v)} a(w)r(\overline{w})
$$
  
\n
$$
= \sum_{x \in T_l} a(x)r(\overline{x})s^{-1}(x) \sum_{v \in S(x)} q(\overline{v}) \sum_{w \in S(v)} q(\overline{w})
$$
  
\n
$$
= \sum_{x \in T_l} a(x)r(\overline{x})s^{-1}(x) \sum_{v \in S(x)} (1 - q(\overline{v})) .
$$

For  $x \in T_l \setminus B_l$  and  $\upsilon \in S(x)$ , (25) and (28) yield  $q(\overline{\upsilon}) \leq d^{-1/2} + d^{-3/4+2\varepsilon}$ . Hence

$$
\sum_{w \in T_{l+2}} a(w)r(\overline{w}) \ge \sum_{x \in T_l \setminus B_l} a(x)r(\overline{x})s^{-1}(x) \sum_{v \in S(x)} (1 - q(\overline{v}))
$$
\n
$$
\ge (1 - d^{-1/2} - d^{-3/4 + 2\varepsilon}) \sum_{x \in T_l \setminus B_l} a(x)r(\overline{x})
$$
\n
$$
\ge (1 - d^{-1/2} - d^{-3/4 + 2\varepsilon}) \sum_{x \in T_l} a(x)r(\overline{x}) - \sum_{x \in B_l} a(x)r(\overline{x})
$$
\n
$$
\ge p(\overline{v}) (1 - d^{-1/2} - d^{-3/4 + 2\varepsilon})^{(l+2)/2} - \frac{l \cdot e^{-d\varepsilon/5}}{2} - \sum_{x \in B_l} a(x)r(\overline{x}).
$$

So it is enough to show that

$$
\sum_{\mathbf{x}\in B_l} a(\mathbf{x})r(\overline{\mathbf{x}}) \leq e^{-d^{\varepsilon}/5}.
$$

However, the Cauchy-Schwarz inequality, (24) and Lemma 5.2 imply that

$$
\sum_{x \in B_l} a(x) r(\overline{x}) \leq \left( \sum_{x \in B_l} a^2(x) \right)^{1/2} \left( \sum_{x \in B_l} r^2(\overline{x}) \right)^{1/2} \leq \left( t(d-t)^l e^{-d^{\varepsilon}/2} (d-t)^{-l} \right)^{1/2} \left( p_G(\overline{v}) \right)^{1/2} \leq e^{-d^{\varepsilon}/5} .
$$

Claim 2.

$$
\sum_{w \in T_l} a^2(w) \le 1 + d^{-1/4 + 2\varepsilon} \quad \text{ for all even } l \le t .
$$

Proof. We show by induction that

$$
\sum_{\mathbf{w}\in T_l} a^2(\mathbf{w}) \le \left(1 + 2d^{-3/4+\varepsilon}\right)^{l/2} \left(1 + \frac{l}{2t}e^{-d^{\varepsilon}}\right). \tag{37}
$$

(As usual, the base case is trivial.) Notice that

$$
\sum_{w \in T_{l+2}} a^2(w) = \sum_{x \in T_l} \sum_{v \in S(x)} \sum_{w \in S(v)} a^2(w)
$$
  
\n
$$
= \sum_{x \in T_l} a^2(x) s^{-2}(x) \sum_{v \in S(x)} \sum_{w \in S(v)} 1
$$
  
\n
$$
= \sum_{x \in T_l} a^2(x) s^{-1}(x) \overline{s}(x)
$$
  
\n
$$
\leq (1 + 2d^{-3/4 + \varepsilon}) \sum_{x \in T_l} a^2(x) + \left(\frac{d}{d - t} - 1\right) \sum_{x \in T_l} a^2(x).
$$

On the other hand, (22) gives

$$
\sum_{x \in \Gamma_l} a^2(x) \le t^{-1} (d-t)^l e^{-d^{\varepsilon}} (d-t)^{-l} = t^{-1} e^{-d^{\varepsilon}},
$$

so we have (37).

 $\Box$ 

*Proof of* (34). Claims 1 and 2 with (36) imply that, for all even  $l \le t/2 - 2$ ,

$$
\sum_{w \in T_l} r^2(\overline{w}) \ge p^2(\overline{v}) \left(1 - d^{-1/2} - 2d^{-3/4 + 2\varepsilon}\right)^l \left(1 + d^{-1/4 + 2\varepsilon}\right)^{-1}
$$
  
 
$$
\ge p^2(\overline{v}) \left(1 - d^{-1/4 + 3\varepsilon}\right) \left(1 - d^{-1/2}\right)^l.
$$

Thus

$$
\sum_{\substack{w \in V(T) \setminus \{v\} \\ |w| \text{ even}}} r^2(\overline{w}) \ge \sum_{l=2}^{t/2-2} \sum_{w \in T_l} r^2(\overline{w})
$$
\n
$$
\ge p^2(\overline{v}) \left(1 - d^{-1/4+3\varepsilon}\right) \sum_{\substack{l=2 \\ l \text{ even}}}^{t/2-2} \left(1 - d^{-1/2}\right)^l
$$
\n
$$
\ge \frac{p^2(\overline{v})(1 - 2d^{-1/4+3\varepsilon})d^{1/2}}{2}
$$
\n
$$
\ge \frac{p(\overline{v})(1 - d^{-1/4+4\varepsilon})}{2}
$$

(where we again use Theorem  $1.2(a)$  in the final inequality).



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